## MATH455 HOMEWORK 4 DUE FRIDAY, FEBRUARY 21

*Exercise* 1. Suppose that a theory T has arbitrarily large finite models. Show that T has an infinite model.

*Exercise* 2. A linear order (I, <) is called a *well-order* if there is no infinite descending sequence  $a_0 > a_1 > \cdots > a_k > \cdots, k \in \mathbb{N}$ , of elements from I.<sup>1</sup>

Show that if  $\mathcal{I} = (I, <)$  is an infinite linear order then there is  $\mathcal{J} \succ \mathcal{I}$  which is *not* a well-order. In particular, this shows that you cannot write down axioms in first-order logic which characterize being a well-order.

*Exercise* 3. Let (I, <) be a linear order and suppose that  $\langle \mathcal{M}_i : i \in I \rangle$  is a sequence of  $\mathcal{L}$ -structures so that if i < j are elements of I then  $\mathcal{M}_i \prec \mathcal{M}_j$ . We call  $\langle \mathcal{M}_i : i \in I \rangle$  an *elementary chain*. Show that if  $\mathcal{M} = \bigcup_{i \in I} \mathcal{M}_i$  then there is a structure  $\mathcal{M}$  with underlying set  $\mathcal{M}$  so that  $\mathcal{M}_i \prec \mathcal{M}$  for all  $i \in I$ . [Hint: first you have to say what the constants, functions, and relations of  $\mathcal{M}$  are.]

*Exercise* 4. Let  $\mathcal{M}$  be an infinite structure. Suppose  $p(x) = \{\varphi(x)\}$  is a set of formulae in the language of  $\mathcal{M}$  who all have x as their only free variable. Suppose that for all finite  $p_0(x) \subseteq p(x)$  there is  $a \in \mathcal{M}$  so that  $\mathcal{M} \models \varphi[a]$  for all  $\varphi(x) \in p_0(x)$ . Show that there is  $\mathcal{N} \supseteq \mathcal{M}$  so that there is  $a \in \mathcal{N}$  with  $\mathcal{N} \models \varphi[a]$  for all  $\varphi(x) \in p(x)$ .

<sup>&</sup>lt;sup>1</sup>For example,  $(\mathbb{N}, <)$  is a well-order but  $(\mathbb{Z}, <)$  and  $(\mathbb{Q}, <)$  are not well-orders.